Parsing & Understanding in a Messy World

Mike Dodds - January 2024

galois

Context: Galois / me

Galois: A contract research shop. Paid research by-the-hour

- Security / reliability technologies (PL, formal methods, static analysis)
- Clients: DARPA, US Gov, some commercial

Me: verified cryptography, formal methods at scale, parser security

- $2004 \rightarrow 2017$: York / Cambridge / York PhD, postdoc, junior professor
- $2017 \rightarrow now$: Galois principal scientist (~ full professor)

Context: DARPA SafeDocs project

SafeDocs: Galois + other teams try to make parsing better

Galois built two tools:

- Format Analysis Workbench (FAW), a tool for understanding existing formats
- Daedalus, a language for developing safer parsers

We have:

- built a high-assurance parser which covers most real-world PDFs
- analyzed millions of real-world PDF documents
- fixed multiple issues in the PDF standard

This talk:

- 1. Parsing matters a lot and is very hard
- 2. Eg: PDF, an interesting and horrible format
- 3. Two core problems in safer parsing
- 4. Some progress: FAW & Daedalus

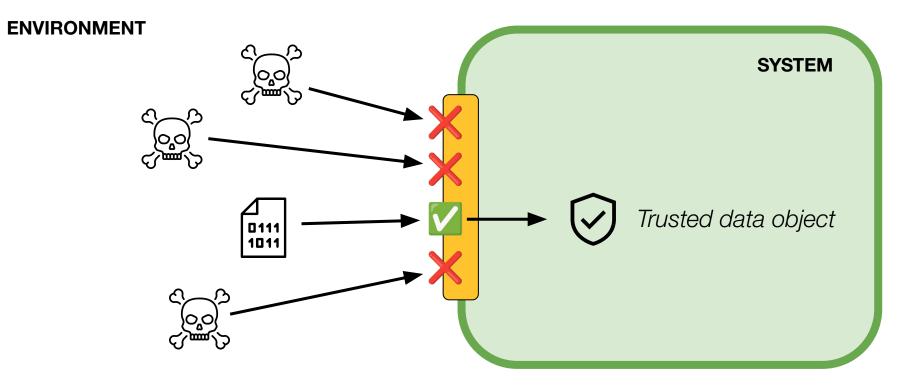
1. Parsing matters a lot and is very hard

Parsers are the immune system

A system has an outside (low trust) and an inside (high trust)

Systems interact with the world

Parsers convert low-trust data to high-trust data

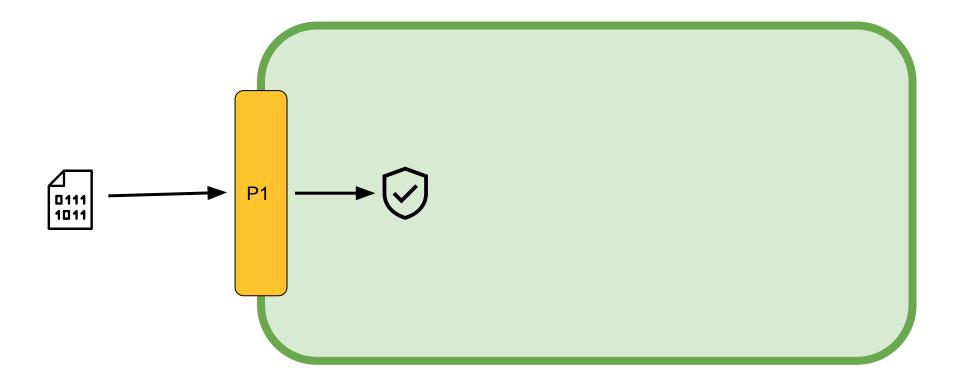


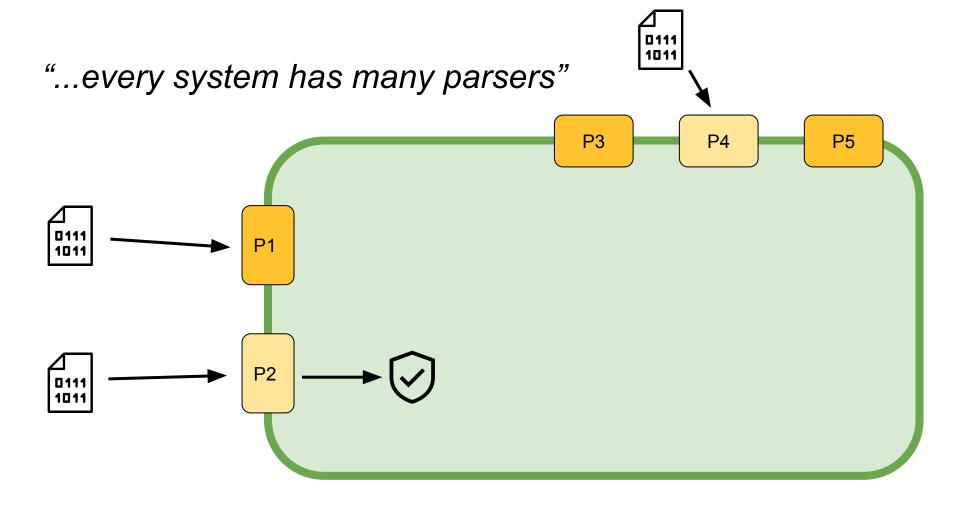
Environment is untrustworthy

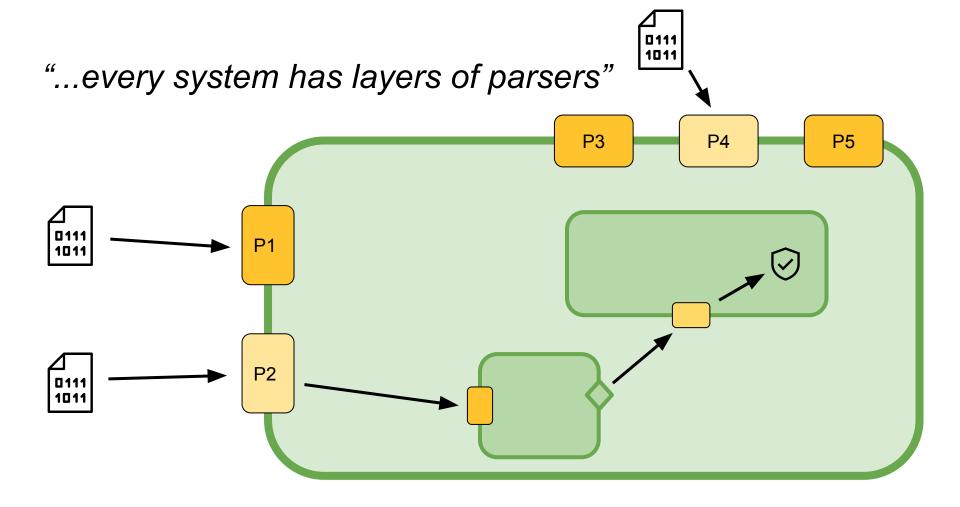
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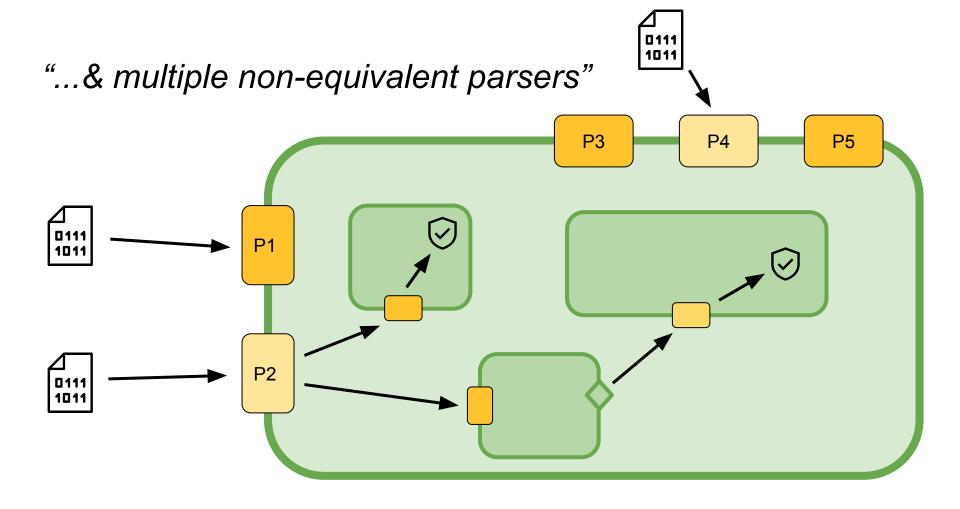
Everyone writes parsers

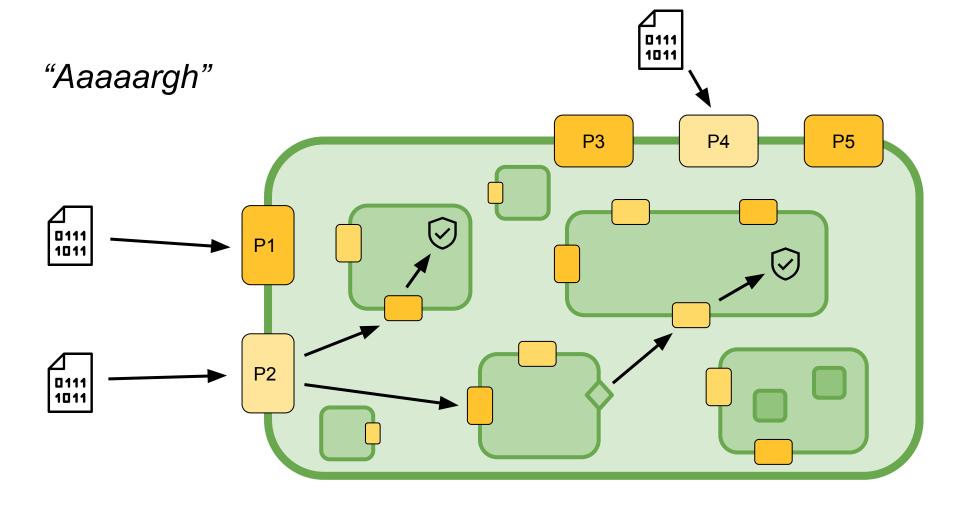
Me (naive): "Every system has a parser"

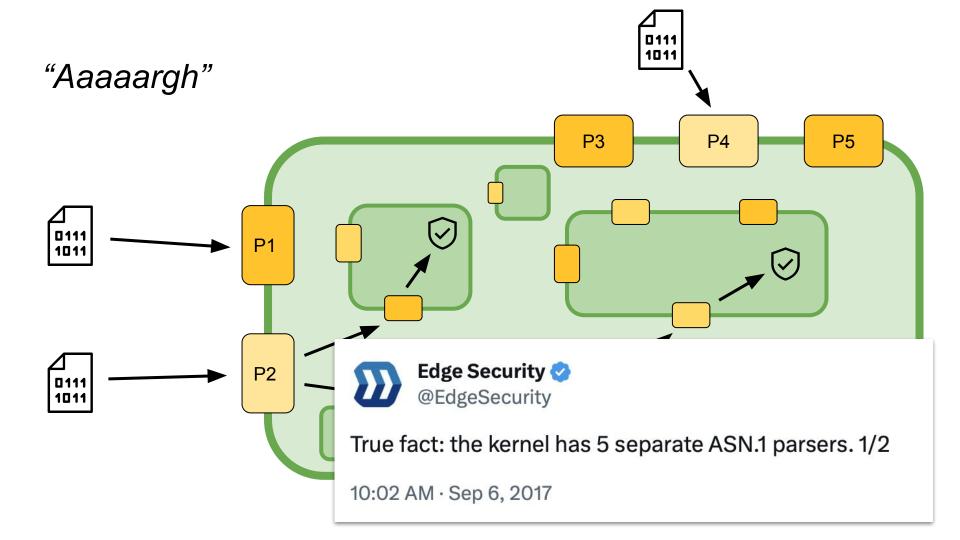












Why are there so many parsers?

Parsers are:

- performance critical
- intermingled with computation
- perform different tasks (security filter, data parsing, constructing values)
- written in languages don't provide clean abstractions

Also: systems are built over time, and parsers tend to grow capabilities

Parsers fail in interesting ways

1: Parsers crash

This is quite bad

Usually this means a memory safety violation

Potentially, this allows an adversary to write into memory

Desired property: absence of undefined behavior

No specification required. A crash is a crash.

2: Parsers construct semantic values incorrectly

This is bad, obviously

Eg. data sanitization failures - see \rightarrow

But also, meaning is reconstructed wrongly

Desired property: the parser behaves as intended

We need to know what is intended

(...even for non-conformant inputs)

DID YOU REALLY NAME YOUR SON Robert'); DROP TABLE Students; -- ? OH, YES. LITTLE BOBBY TABLES, WE CALL HIM.

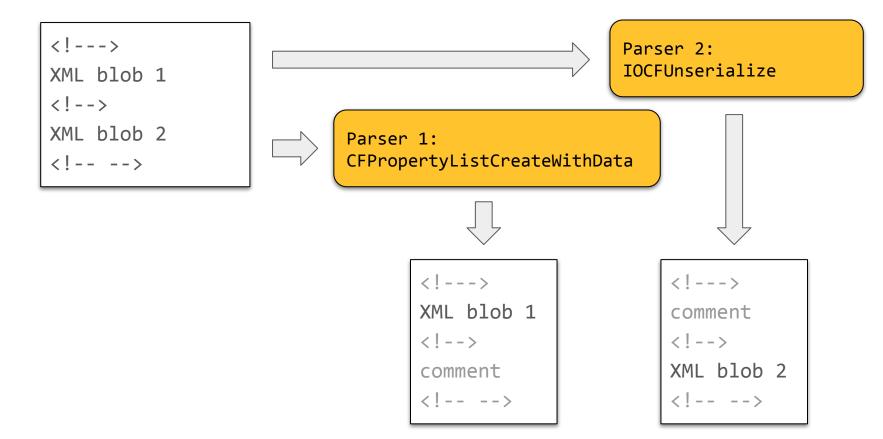
https://xkcd.com/327/

3: Parsers disagree

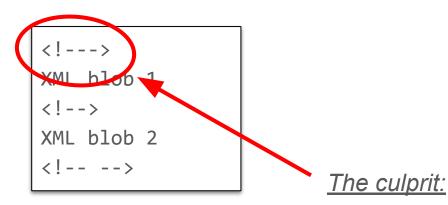
Er, this might be bad?

Actually, let me give an example...

Aside: 'psychic paper' https://blog.siguza.net/psychicpaper/



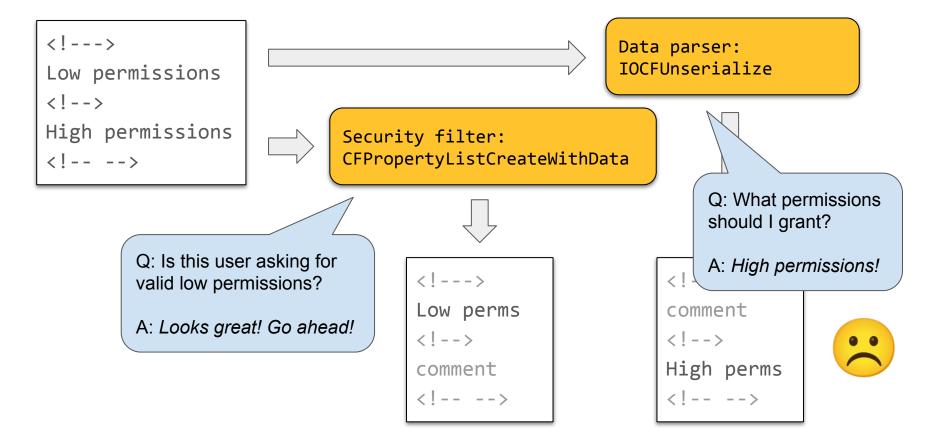
Aside: 'psychic paper' https://blog.siguza.net/psychicpaper/



Not valid XML. It could be interpreted as:

- "Start comment"
- "Start and end comment"
 - \Rightarrow Both parsers are 'correct'!

Aside: 'psychic paper' https://blog.siguza.net/psychicpaper/



3: Parsers disagree

So this is actually bad, and hard to detect

Examples:

- Sneak past security parsers
- PDFs that parse differently when viewed and printed

Desired property: parsers agree with each other

This is a meta-property between parsers

Writing correct parsers is very hard

Ground truth does not exist

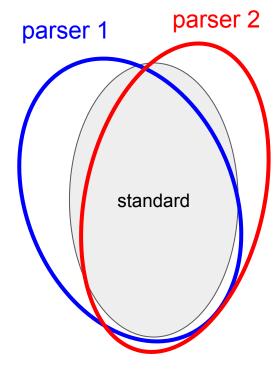
We can assume:

- One or more existing parsers
- Some documentation, and maybe a standard
- A set of examples of the format

We can't assume:

- Agreement between existing parsers
- Specifications matching *de facto* behavior

Parsers are incentivized to parse non-conformant inputs



Common parser languages are difficult to audit

Typically, parsers are written in C++ & similar:

- Hard to even establish absence of undefined behavior
- Hard to extract parser behavior / reason about parsers
- Hard to specify parsers at the high level
- Hard to audit behavior for humans

Fuzzing is good but limited

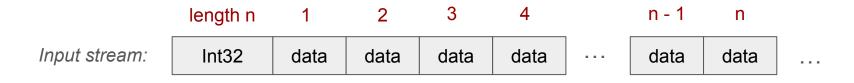
- Generate lots of random inputs
- Guided search for crashes

Fuzzing is de-facto the way that parsers are secured

But: only finds *undefined behavior* parser flaws

Hold on didn't we solve parsing in, like... 1959?

Damn you, data dependency



We have to:

- Parse length, and compute n
- Read **n** more data chunks

In general, parsing has to perform arbitrary computations

I'll come back to this later...

"Shotgun parsing" (Brattus et al)

A common parser structure for dealing with data dependency:

- Read some data
- Call some arbitrary handler function (e.g written in C++)
- Return a value and keep going

Unsafe, hard to maintain, non-auditable

2. PDF, an interesting and horrible format

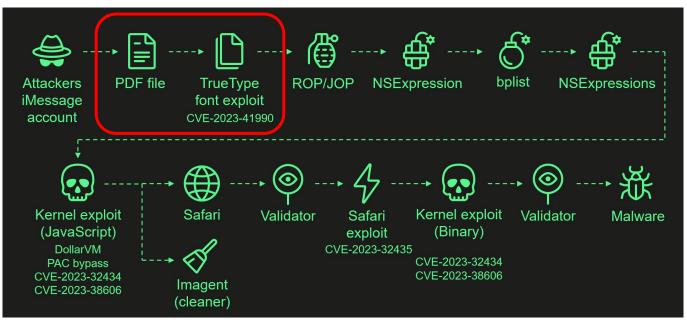
PDF is important and interesting

- Billions of users
- *De facto* message format for many human processes
- Huge attack surface & many vulnerabilities
- Contains embedded formats: images, fonts, JavaScript, video (...yes, really)

- Has a somewhat agreed core standard
- Many real implementations some good, some v bad
- Huge dataset of examples in the wild

PDF is an attack vector

Eg. Operation Triangulation (December 2023) https://securelist.com/trng-2023/



Operation Triangulation exploit chain

SafeDocs built a huge dataset of PDFs

https://pdfa.org/new-large-scale-pdf-corpus-now-publicly-available/

About 8m extant (real-world) documents

A large proportion of these PDFs don't conform to the standard

PDF is weird and hard to parse

XRef table

PDF structure:

- A set of objects text, values, pages, fonts...
- A cross-reference table of object locations

XRef supports incremental updates

Surprising results:

- Parsing is highly non-linear
- Parsing depends on computing offsets

Header

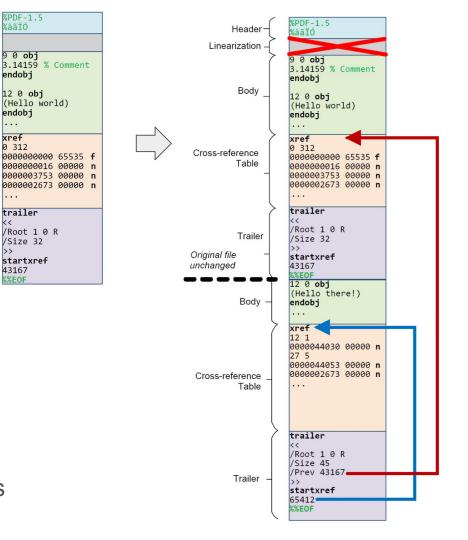
Body

Table

Trailer

Linearization

Cross-reference



Object streams

Objects can be contained in other objects

- Object sizes can be contained in other objects
- Objects can be compressed or encrypted

Surprising results:

- Parsing an object may require accessing multiple other objects
- Parsing may require decrypting / decompressing other object

```
9 0 obj
3.14159 % Comment
endobj
12 0 obj
(Hello world)
endobj
...
```

Themes in parsing PDF:

Pervasive data-dependency Pervasive computation Non-local parsing Many embedded formats

Hypothesis: most mature formats are super weird

3. Two core problems in safer parsing

"What do existing parsers do?"

"How can we write better parsers?" Implement parser understanding

"What do existing parsers do?"

"How can we write better parsers?" Implement parser understanding

"What do existing parsers do?"

"How can we write better parsers?"

Test new safer parsers

Some progress (Daedalus and FAW)

"What do existing parsers do?"

"How can we write better parsers?"

Format Analysis Workbench

An *investigation engine* for understanding parsers and formats

FAW can:

- Run parsers at scale
- Analyze results
- Test hypotheses
- Generate understanding

Daedalus

A *format description language* for generating safe and correct parsers

Daedalus can:

- Define human-readable format definitions
- Prevent crashes
- Synthesize parsers

Format Analysis Workbench

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Daedalus

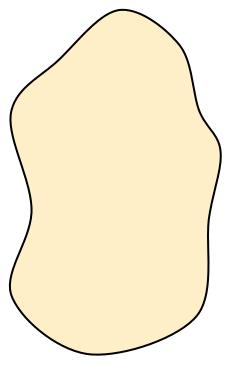
A *format description language* for generating safe and correct parsers

Daedalus can:

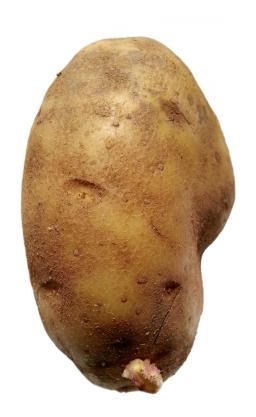
- Define human-readable format definitions
- Prevent crashes
- Synthesize parsers

Format Analysis Workbench (FAW)

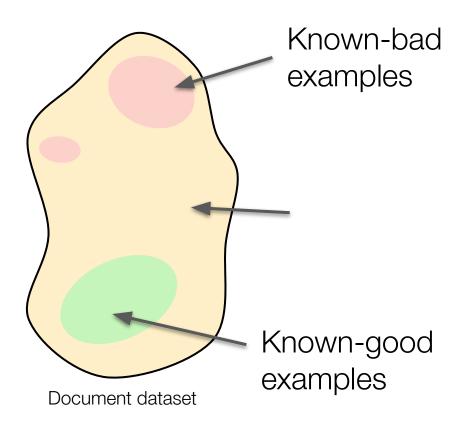
"What do existing parsers do?"



Document dataset

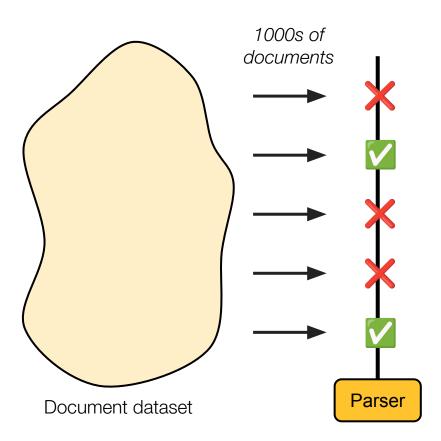


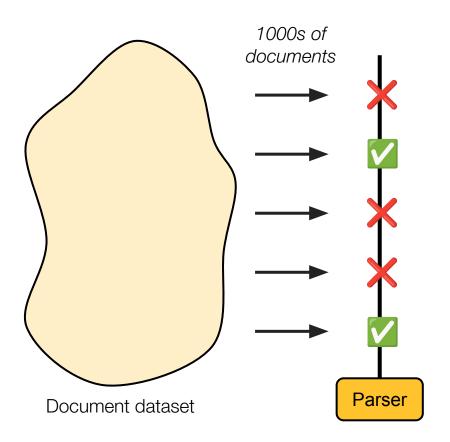
Document dataset



Eg, PDF dataset:

- 1M+ files
- Some known-good and known-bad examples, but mostly unknown



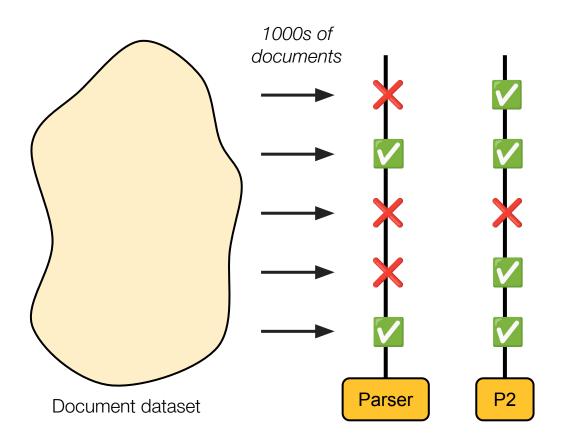


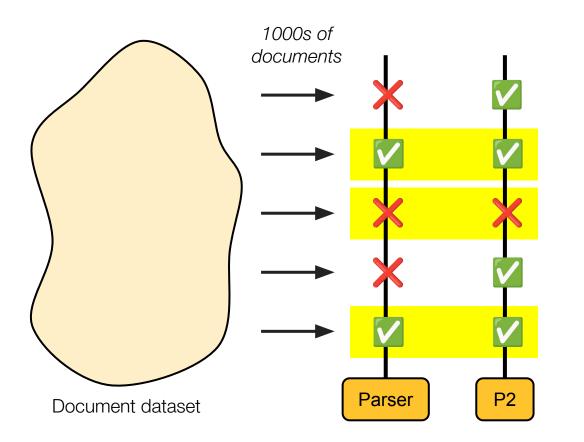
Results of parsing:

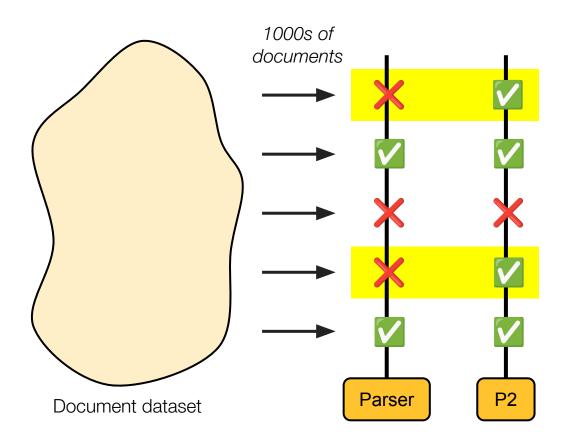
- Valid 🔽 or invalid 🗙
- Parser return codes

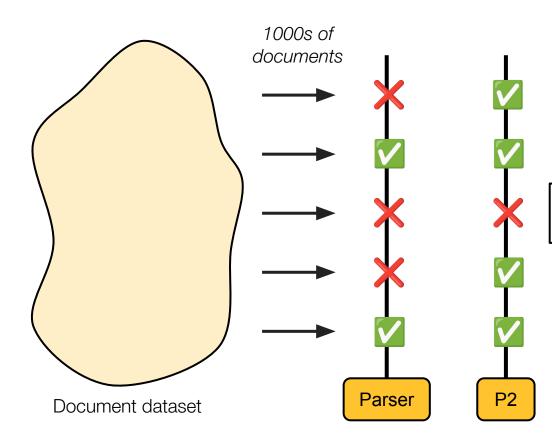
But also:

- Plug-in analysis results
- Any tool that can apply to a parser!









Parser \\ P2	rejected	valid
rejected	X %	A %
valid	В%	Υ%

The FAW is a format science lab

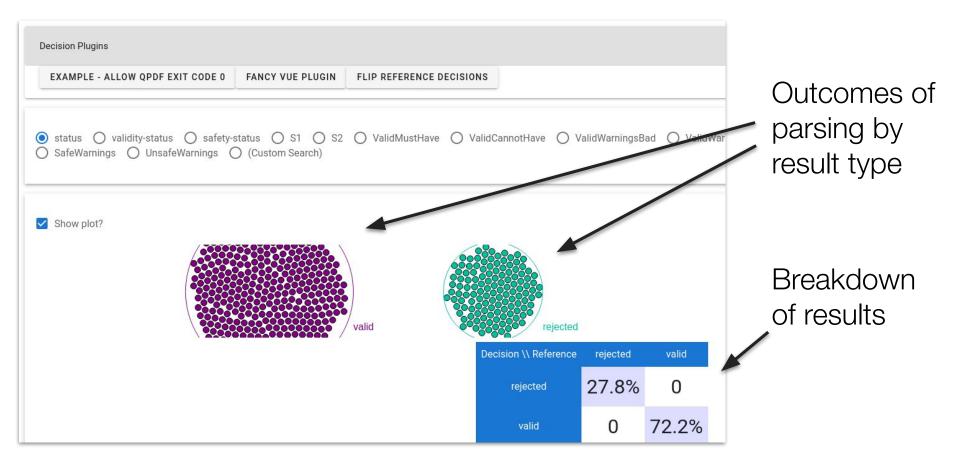
Inputs:

- Format examples (e.g., PDFs)
- Parsers or programs that ingest those examples

Use cases:

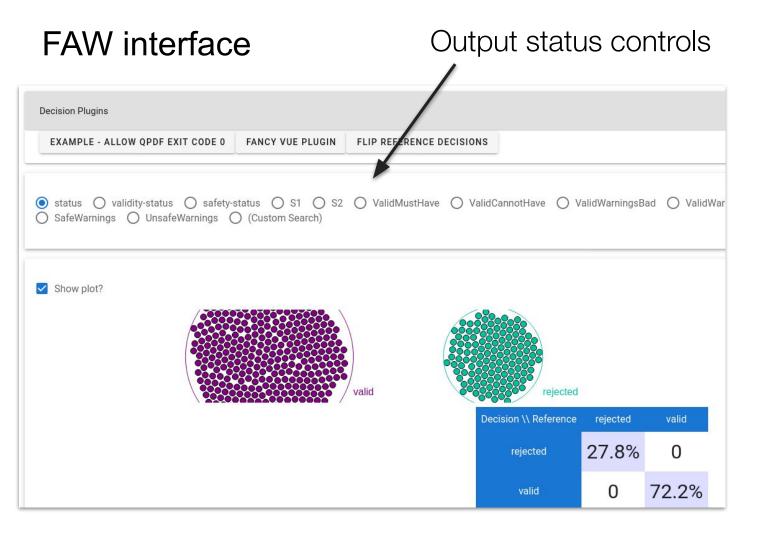
- Identify potentially unsafe inputs
- Identify causes of false alarms at scale
- Understand patterns of input and how they affect individual programs at a deep level

FAW interface

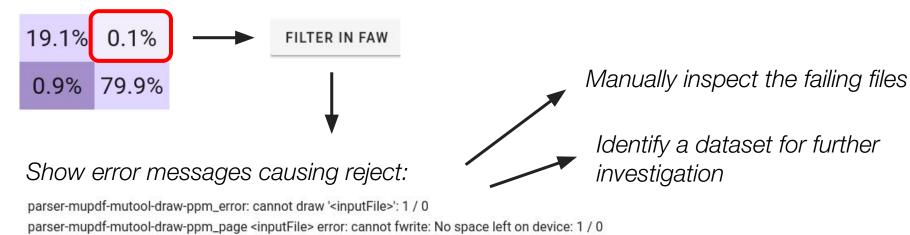


Categorization is controlled by the user

```
outputs:
 # Standard output status -- If a PDF passes filter S1, it will be "valid",
 # otherwise "rejected".
 status:
   "valid" is !(RejectedBad | RejectedAmbiquousBad | ValidWarningsXrefRebuild & XrefIsAmbiquous)
   "rejected" else
 # Mark anything that is linearized as "rejected-unsafe", and otherwise
 # mark it as "valid".
 validity-status:
   "rejected" is RejectedBad
   "rejected-ambiguous" is RejectedAmbiguousBad | (ValidWarningsXrefRebuild & XrefIsAmbiguous)
   "valid-warnings" is ValidWarningsBad | !ValidMustHave | ValidCannotHave | ValidWarningsXrefRebuild
   "valid" else
   "rejected-unsafe"
  safety-status:
                                                Output status is determined by
   "unsafe"
   "unsafe-warnings" is UnsafeWarnings
   "safe-warnings" is SafeWarnings
                                                a regexp-based alarm language
   "safe" else
```



The FAW assists in format detective work



parser-xpdf-pdftoppm_I/O Error: Couldn't create temporary font file: 1 / 1

parser-mupdf-mutool-clean_error: cannot fwrite: No space left on device: 3 / 3

Run further analyses and discover correlations

Interrogation example: PolyFile

A utility by *Trail of Bits* for examining the structure of files and detecting their file type

- Hex viewer for examining the file in detail which shows how various parts of the binary are interpreted
- Map file data back to AST nodes generated by the parser
- Plugs into FAW

https://github.com/trailofbits/polyfile

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Interrogation example: PolyTracker

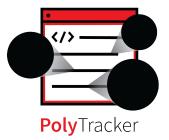
https://github.com/trailofbits/polytracker

PolyTracker: a generic taint tracking tool

- Binary instrumentation to track data processing
- Identify which parser functions touch which parts of the file

PolyFile + PolyTracker:

- Map a file's meaning (PolyFile)
- Map how it is used in the binary (PolyTracker)



"What do existing parsers do?"

Implement parser understanding

"What do existing parsers do?"

"How can we write better parsers?"

"How can we write better parsers?"

Daedalus

Daedalus: a language for writing formats

Aim: close the gap from formats to parsers

- Powerful enough to represent eg. PDF
- Amenable to human reading and static analysis
- Type-safe, crashes can't happen
- Turing-complete, but highly structured
- Amenable to performant compilation into C++

Daedalus is a language and toolchain

Daedalus (language): Data Description Language

Daedalus (toolchain): compilation and execution of Daedalus-lang specifications

• Compile Daedalus to performant C++ code

Daedalus design

Based on functional programming ideas / parser combinators

Includes several highly useful capabilities:

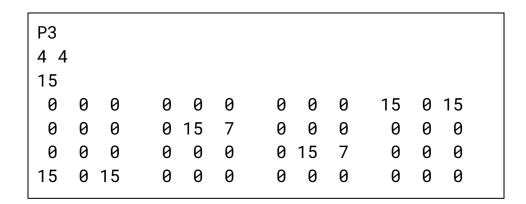
- A generic notion of data dependency. Depend on any datatype
- An encapsulated notion of an input stream. Safe non-linear parsing.
- An FFI interface. Call into helper functions in a controlled way

Example: PPM, a small image format

Specification:

- A magic number identifying the file type (for ASCII PPM, this is *P3*)
- The dimensions of the image (width then height)
- The maximum color value
- A 'matrix' of RGB triples for each pixel defined in row-major order

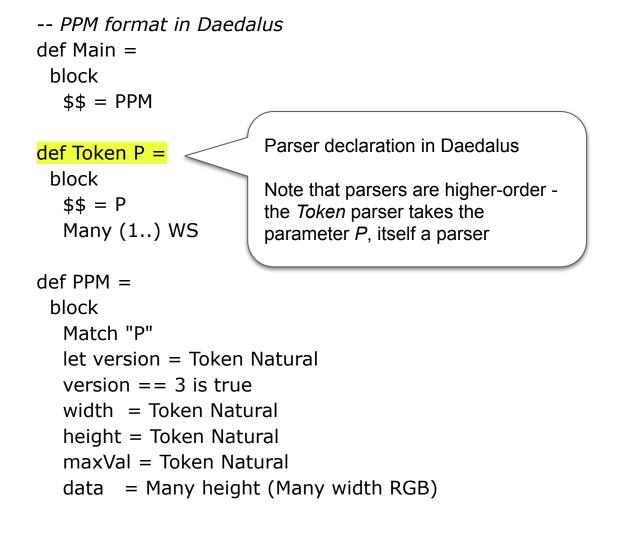
A PPM file



- The magic number is P3, indicating an ASCII RGB image
- The width and height are both 4
- The maximum color value is 15
- There is a four-by-four grid of triples, one triple per pixel

-- PPM format in Daedalus def Main =block \$\$ = PPMdef Token P =block \$\$ = P Many (1..) WS def PPM =block Match "P" let version = Token Natural version == 3 is true width = Token Natural height = Token Natural maxVal = Token Natural data = Many height (Many width RGB)

def RGB = block red = Token Natural green = Token Natural blue = Token Natural def WS = Match1 (0 | 9 | 12 | 32 | 'n' | 'r') def Natural = block let ds = Many (1..) Digit $^{\text{for}}$ (val = 0; d in ds) (addDigit val d) def Digit = block let d = Match1 ('0' ... '9')'0' - b ^



PPM format in Daed def Main = block \$\$ = PPM	dalus
def Token P = block \$\$ = P Many (1) WS	
def PPM =	
block	Primitive parsing in Daedalus
Match "P" let version = Toke version == 3 is tr	The parser reads a token "P" off the input stream
width = Token Na height = Token Na maxVal = Token Na	If no such token is present, the parser backtracks
data 🛛 = Many heig	ht (Many width RGB)

Primitive parsing with multiple possible values

The parser *WS* reads one of the possible choices: 0, 9, ...

```
def RGB =
```

```
block
```

red = Token Natural

green = Token Natural

blue = Token Natural

def WS = Match1 (0 | 9 | 12 | 32 | '\n' | '\r')

```
def Natural =
  block
  let ds = Many (1..) Digit
  ^ for (val = 0; d in ds) (addDigit val d)
```

```
def Digit =
   block
   let d = Match1 ('0' .. '9')
    ^ d - '0'
```

Parser combinators in Daedalus

The *red, green,* and *blue* values are parsed in sequence using the *block* combinator

The return type of the block is a *structure type* with fields *red, green, blue*

```
def RGB =
block
```

```
red = Token Natural
green = Token Natural
blue = Token Natural
```

def WS = Match1 (0 | 9 | 12 | 32 | '\n' | '\r')

def Natural =
 block
 let ds = Many (1..) Digit
 ^ for (val = 0; d in ds) (addDigit val d)

```
def Digit =
   block
   let d = Match1 ('0' .. '9')
    ^ d - '0'
```

def RGB =

block

red = Token Natural

green = Token Natural

blue = Token Natural

def WS = Match1 (0 | 9 | 12 | 32 | '\n' | '\r')

Computation in Daedalus The *Natural* parser reads multiple digits, and then computes the overall value by iterating over the

list of digits

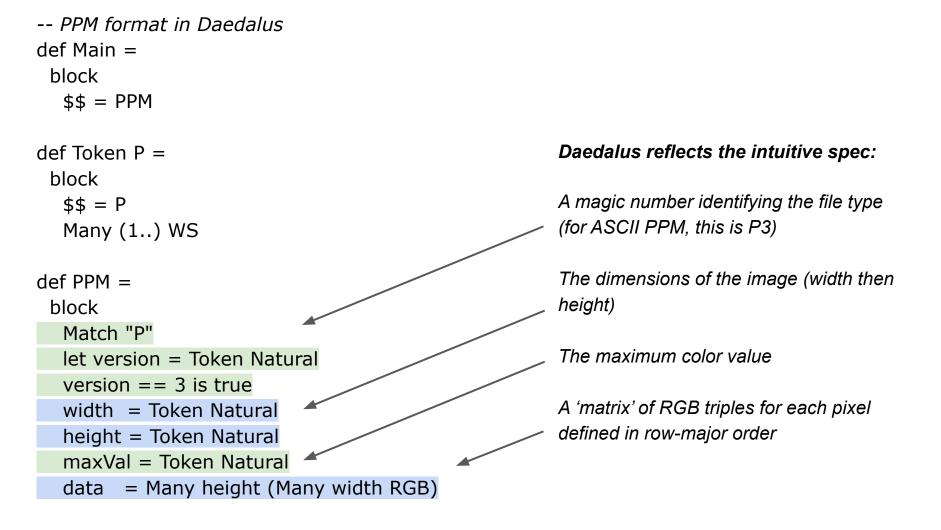
def Natural =
 block
 let ds = Many (1..) Digit
 ^ for (val = 0; d in ds) (addDigit val d)

```
def Digit =
   block
   let d = Match1 ('0' .. '9')
   ^ d - '0'
```

-- PPM format in Daedalus def Main = block \$\$ = PPMdef Token P =block \$\$ = P Many (1..) WS def PPM =block Match "P" let version = Token Natural version == 3 is true width = Token Natural height = Token Natural maxVal = Token Natural data = Many height (Many width RGB)

Data dependency in Daedalus

The parser behaviour depends on the *width* and *height* values computed during earlier parsing





Tutorial: <u>https://galoisinc.github.io/daedalus/tutorial/index.html</u>

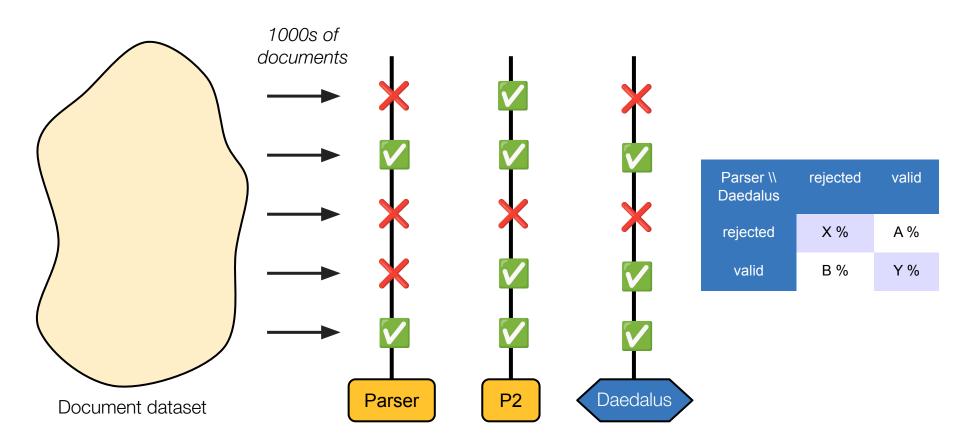
try-Daedalus, a framework for developing Daedalus in VSCode, using a remote container: <u>https://github.com/galoisinc/try-Daedalus</u>

FAW + Daedalus

Implement parser understanding

Format Analysis Workbench (FAW) Daedalus language & toolchain

Test new safer parsers



We tested FAW + Daedalus a lot!

- Daedalus definitions and generated parsers for **14 formats (inc PDF)**
- Analyzed **13 PDF parsers**, **5 NITF parsers** and **1MM+ documents**
- Discovered 9 issues with PDF specification, 10s of bugs in parsers
- Working with the PDF Foundation to develop a machine-readable specification of PDF that eliminates common vulnerabilities

We built other things (thanks, DARPA!)

Talos, an object synthesizer based on symbolic analysis of Daedalus specs

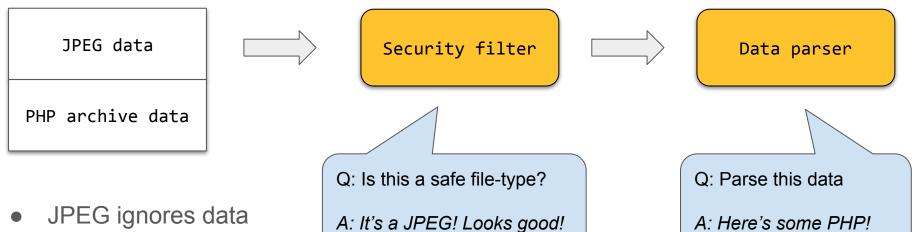
- Daedalus: parse bytes into a semantic value
- Talos: from a format and desired semantic value, construct the input bytes

HTTP smuggling detection, FAW + Talos to find HTTP parser differentials

Polyglot detection, based on Daedalus + static analysis

Polyglots are bad

file:



9 9

- JPEG ignores data after end
- PHP archive ignores data before 'magic' start string

Polyglots are caused by *cavities* (& other things too)

AB type: <cavity> Data

Zipper type:	
<cavity></cavity>	
Data	
<cavity></cavity>	
Data	

- Cavities don't affect the resulting semantic value
- Caused by eg. comments, start characters

Eg: Evan Sultanik @ Trail of Bits - resume is PDF and NES ROM:

https://www.sultanik.com/cv

Static cavity detection in formats

Cavity detection process:

- Write the format in Daedalus
- Use a context-sensitive, flow-insensitive analysis to track how parsed data is handled
- Cavities form when data is parsed but not tested by the program

Detect potential polyglots based on a Daedalus format description

Also: synthesize polyglot instances using Talos

Daedalus as a target for format analysis!

Summary:

- 1. Parsing matters a lot and is very hard
- 2. Eg: PDF, an interesting and horrible format
- 3. Two core problems in safer parsing
- 4. Some progress: FAW & Daedalus

SafeDocs team

- Galois (Prime)
- Trail of Bits
- RTI
- **Narf Industries**
- Verocel
- Cornell
- Penn State
- Princeton
- Purdue
- Tufts





Parsing is, unfortunately, still very hard

Some problems we thought about but didn't solve:

- Parser verification, especially for extant parsers
- Subsetting / filtering parser languages
- Managing variants of formats (eg. a spec vs non-conformant versions)
- Specifying schema descriptions (eg. JSON) alongside data formats
- Synthesis of format specifications from examples

Someone should solve these problems too



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